

Relational Model

Exercises

3.1 Answer: One solution is described below. Many alternatives are possible. Underlined attributes indicate the primary key.

student (student-id, name, program)
 course (courseno, title, syllabus, credits)
 course-offering (courseno, secno, year, semester, time, room)
 instructor (instructor-id, name, dept, title)
 enrolls (student-id, courseno, secno, semester, year, grade)
 teaches (courseno, secno, semester, year, instructor-id)
 requires (maincourse, prerequisite)

3.5 Answer:

- e. $\Pi_{person-name} ((employee \bowtie manages)$
 $\bowtie_{(manager-name = employee2.person-name \wedge employee.street = employee2.street$
 $\wedge employee.city = employee2.city)} (\rho_{employee2}(employee)))$
- f. The following solutions assume that all people work for exactly one company. If one allows people to appear in the database (e.g. in *employee*) but not appear in *works*, the problem is more complicated. We give solutions for this more realistic case later.
- $\Pi_{person-name} (\sigma_{company-name \neq \text{"First Bank Corporation"}}(works))$
 If people may not work for any company:
 $\Pi_{person-name}(employee) - \Pi_{person-name}$
 $(\sigma_{(company-name = \text{"First Bank Corporation"})}(works))$
- g. $\Pi_{person-name} (works) - (\Pi_{works.person-name} (works$
 $\bowtie_{(works.salary \leq works2.salary \wedge works2.company-name = \text{"Small Bank Corporation"})}$
 $\rho_{works2}(works)))$

3.7 Answer:

- a. The left outer theta join of $r(R)$ and $s(S)$ ($r \bowtie_{\theta} s$) can be defined as
 $(r \bowtie_{\theta} s) \cup ((r - \Pi_R(r \bowtie_{\theta} s)) \times (null, null, \dots, null))$
 The tuple of nulls is of size equal to the number of attributes in S .

3.8 Answer:

- c. The update syntax allows reference to a single relation only. Since this update requires access to both the relation to be updated (*works*) and the *manages* relation, we must use several steps. First we identify the tuples of *works* to be updated and store them in a temporary relation (t_1). Then we create a temporary relation containing the new tuples (t_2). Finally, we delete the tuples in t_1 , from *works* and insert the tuples of t_2 .

$$\begin{aligned} t_1 &\leftarrow \Pi_{works.person-name,company-name,salary} \\ &\quad (\sigma_{works.person-name=manager-name}(works \times manages)) \\ t_2 &\leftarrow \Pi_{person-name,company-name,1.1*salary}(t_1) \\ works &\leftarrow (works - t_1) \cup t_2 \end{aligned}$$

3.13 Answer:

- a. $\{t \mid \exists q \in r (q[A] = t[A])\}$
 b. $\{t \mid t \in r \wedge t[B] = 17\}$
 c. $\{t \mid \exists p \in r \exists q \in s (t[A] = p[A] \wedge t[B] = p[B] \wedge t[C] = p[C] \wedge t[D] = q[D] \wedge t[E] = q[E] \wedge t[F] = q[F])\}$
 d. $\{t \mid \exists p \in r \exists q \in s (t[A] = p[A] \wedge t[F] = q[F] \wedge p[C] = q[D])\}$

3.14 Answer:

- a. $\{ \langle t \rangle \mid \exists p, q (\langle t, p, q \rangle \in r_1) \}$
 b. $\{ \langle a, b, c \rangle \mid \langle a, b, c \rangle \in r_1 \wedge b = 17 \}$
 c. $\{ \langle a, b, c \rangle \mid \langle a, b, c \rangle \in r_1 \vee \langle a, b, c \rangle \in r_2 \}$
 d. $\{ \langle a, b, c \rangle \mid \langle a, b, c \rangle \in r_1 \wedge \langle a, b, c \rangle \in r_2 \}$
 e. $\{ \langle a, b, c \rangle \mid \langle a, b, c \rangle \in r_1 \wedge \langle a, b, c \rangle \notin r_2 \}$
 f. $\{ \langle a, b, c \rangle \mid \exists p, q (\langle a, b, p \rangle \in r_1 \wedge \langle q, b, c \rangle \in r_2) \}$

3.19 Answer: To insert the tuple (“Johnson”, 1900) into the view *loan-info*, we can do the following:-

$$borrower \leftarrow (“Johnson”, \perp_k) \cup borrower$$

$$loan \leftarrow (\perp_k, \perp, 1900) \cup loan$$

such that \perp_k is a new marked null not already existing in the database.

Note: no commercial database system supports marked nulls, but strings (or other values) that cannot possibly occur as real values can be used to simulate marked nulls.